

Notation

$X_i \in [0, 1]^d$, $X^N = \{X_1, \dots, X_N\}$, $X^\infty = \{X_1, X_2, \dots\}$,
 \mathcal{R} = Set of axis-aligned hyper-rectangles in $[0, 1]^d$,
 $\text{vol}(\cdot)$ = Lebesgue Measure.

Quantity of interest: Rate of the discrepancy sequence

$$\text{Dis}(X^N) = \sup_{R \in \mathcal{R}} \left| \frac{1}{N} |R \cap X^N| - \text{vol}(R) \right|.$$

Set-up and Objectives

Goal: Design a strategy, which *selects a subset* Z^∞ from a streaming sequence X^∞ of i.i.d. $U[0, 1]^d$ random variables such that $\text{Dis}(Z^N) \ll \text{Dis}(X^N)$.

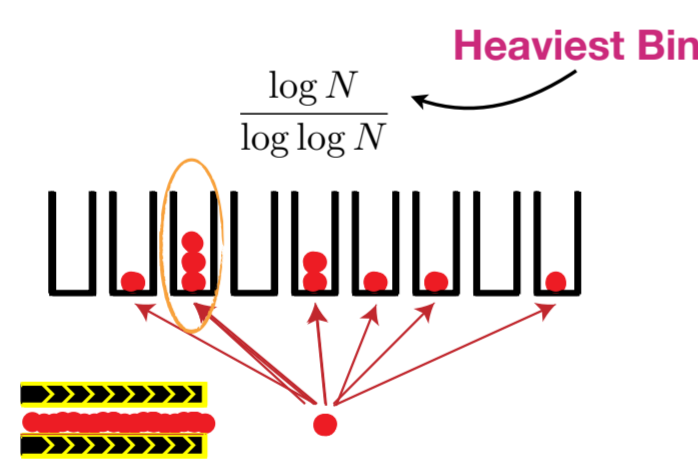
Requirements: We can keep or reject any point in X^∞ . The sequence Z^∞ should be **dense** in X^∞ . The strategy should be **online** and **time and space efficient**.

Remark: $\text{Dis}(X^N) = \text{Monte Carlo Discrepancy} = \mathcal{O}\left(\frac{1}{\sqrt{N}}\right)$.

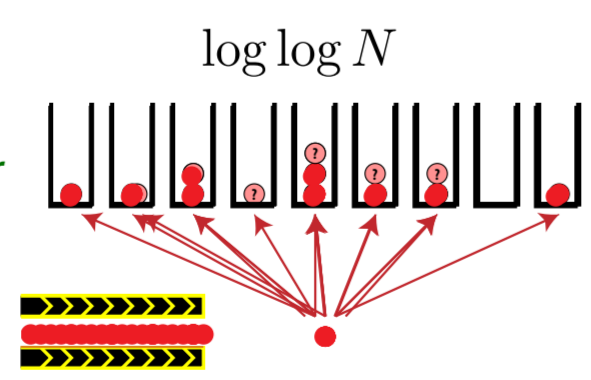
Inspiration: Power of Two Choice

N Balls, N Bins

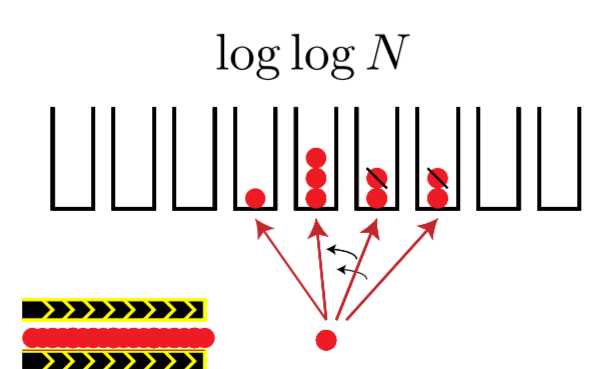
Random Assignment
Randomly sample a bin. Assign the new ball to it.



Two Choice Assignment
Randomly sample two bins. Assign the new ball to the *better* one.



Two Thinning Assignment
Randomly sample a bin. *With some probability* assign the new ball to it. Else, assign the ball to the next random bin.



Idea: Extend the two thinning strategy to our set-up.

Difficulty: The cardinality of \mathcal{R} and heavy dependence across its elements.

Tool: Use of Haar Wavelet Basis.

Main Result

For a streaming sequence X^∞ of i.i.d. $U[0, 1]^d$ random variables, *Haar 2-Thinning strategy* outputs a streaming sequence Z^∞ such that

$$\text{Dis}(Z^N) = \mathcal{O}\left(\frac{d \log^{2d+1}(N)}{N}\right) \quad \forall N \in \mathbb{N}, \text{ almost surely.}$$

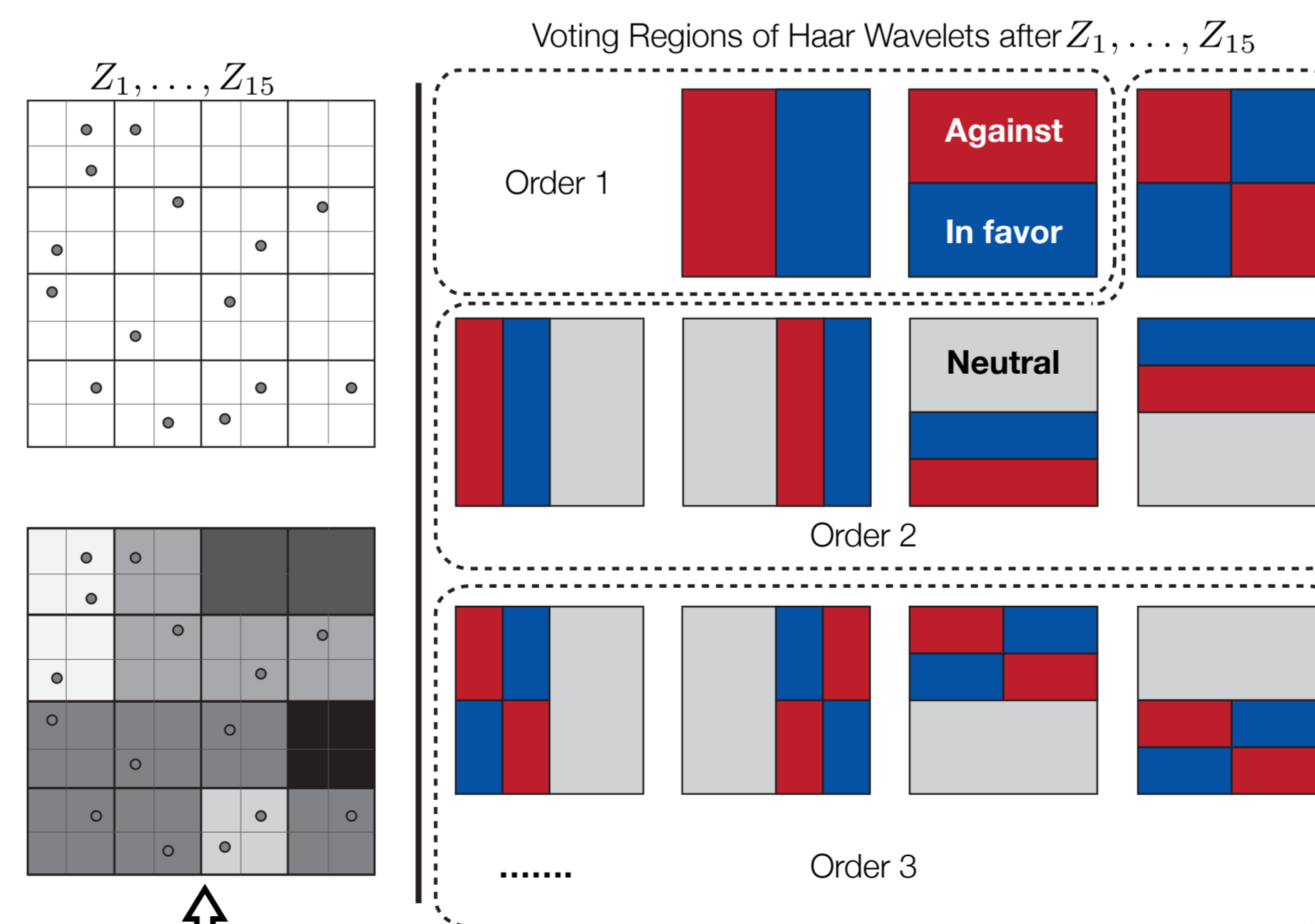
The strategy has $\mathcal{O}(d \log^d N)$ space-time complexity to output N points and **keeps at least one out of every two consecutive points** of the sequence X^∞ .

Haar 2-Thinning Strategy

- Given Z^N and a new candidate point x , each Haar wavelet of order up-to $\lceil \log_2 N \rceil$ votes whether to keep x or reject it.
- Haar functions vote to maintain *balance of points* in their support.
- Based on votes, the point x is kept *randomly with probability*

$$\mathbb{P}(\text{keep } x) = \frac{1}{2} + \frac{1}{2} \cdot \frac{\#\text{votes in favor of } x - \#\text{votes against } x}{\#\text{total votes}};$$

if x is rejected, the next candidate point is kept.



Darker Shade at x = Higher $\mathbb{P}(\text{keep } x)$

Figure 1: Illustration of Haar 2-thinning strategy after obtaining 15 samples in two dimensions. For clarity, $\mathbb{P}(\text{keep } x)$ is averaged on diadic squares of side $1/4$. The strategy favors to keep points in regions that are *deficient* in samples so far.

Haar Greedy-Thinning Strategy

Conjecture: A simplified greedy strategy with $\mathcal{O}(d \log^d N)$ space-time complexity outputs a thinned sequence Z^∞ from a streaming sequence of i.i.d. $U[0, 1]^d$ random variables, such that

$$\text{Dis}(Z^N) = \mathcal{O}\left(\frac{d \log^d N}{N}\right) \quad \forall N \in \mathbb{N}, \text{ almost surely.}$$

Greedy Strategy: Winner takes it all. Keep x *deterministically* if $\#$ votes in favor of $x > \#$ votes against x . Reject otherwise.

Numerical Experiments

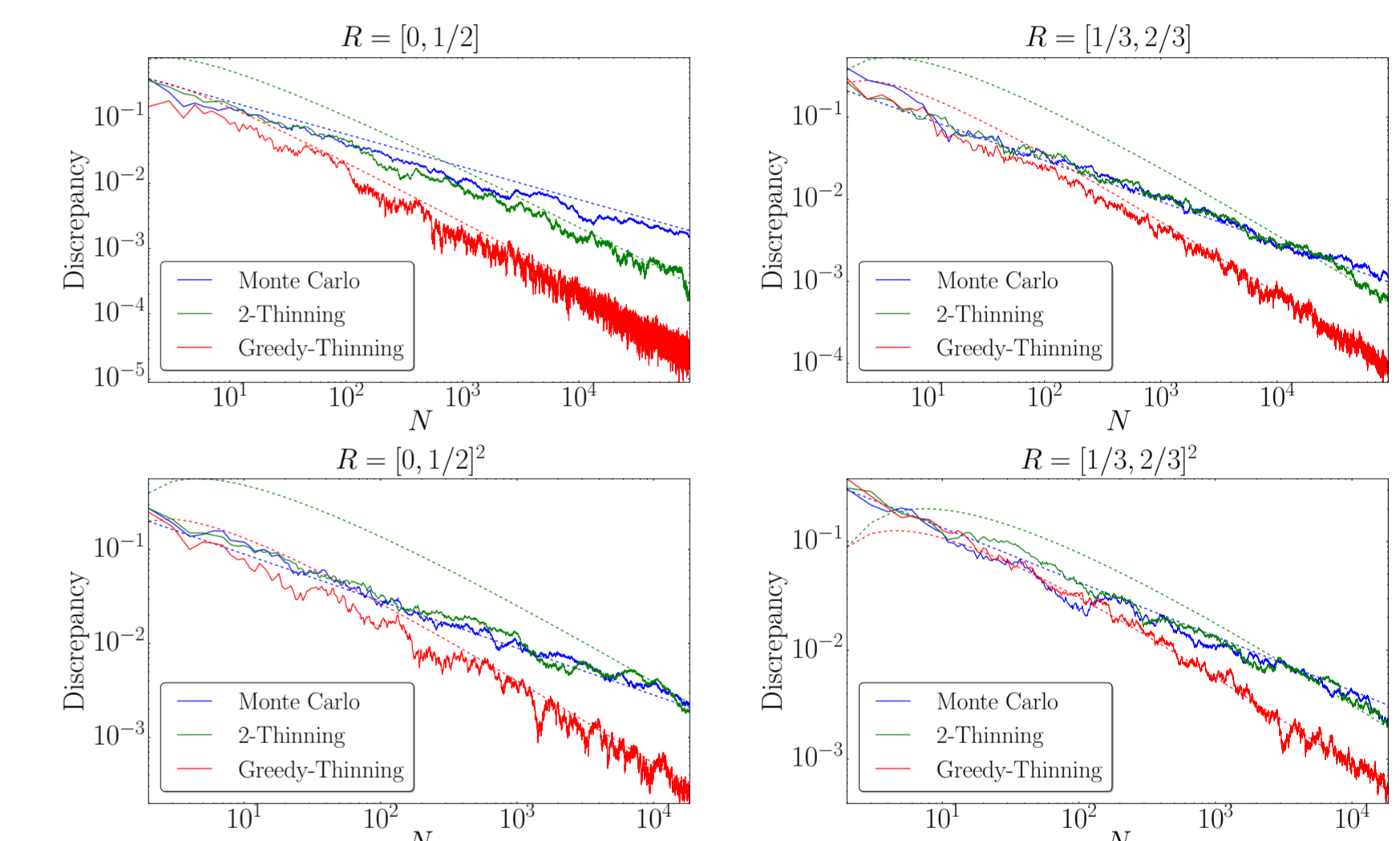


Figure 2: Plots of discrepancy for two different rectangles R in one and two dimensions (averaged over 20 experiments). Axes are in logarithmic scale.

Proof Techniques

- Decompose arbitrary rectangles in terms of “diadic rectangles”
- Express diadic rectangles in terms of “Haar wavelets”
- Exploit *exponential concentration* of self-regulating processes to maintain the balance of points in Haar wavelets up-to some resolution

References

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